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# Modern Logistics & Supply Chain Management

## ML & SCM

*As gold which he cannot spend  
will make no man rich,  
so knowledge which he cannot apply  
will make no man wise.*

Samuel Johnson: The Idler No. 84

*Shortest path problems  
Lie at the heart of network flows.*

Ahuja, Magnanti, Orlin – Network Flows

Shortest Route and  
Maximal Flows

Dr. Wolfgang Garn  
Winter 2016

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# Learning Objectives

- To be familiar with a selected number of classical and state-of-the-art methods
- To create solution models
- To develop procedures that offer competitive advantage to the business/organisation

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# Contents

❖ Shortest path

- An example
- Algorithm
- Integer Program
- Applications

❖ Maximal Flow

- An example
- Algorithm
- Integer Program
- Applications



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### The Shortest Route Problem

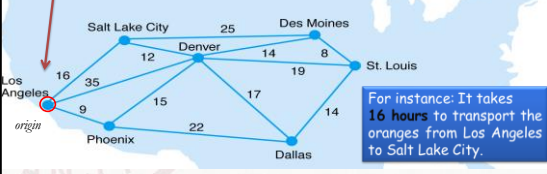
#### Definition and Example Problem Data (1 of 2)

Problem: Determine the shortest routes from the origin to all destinations.

Oranges

6 lorries

6 cities



For instance: It takes 16 hours to transport the oranges from Los Angeles to Salt Lake City.

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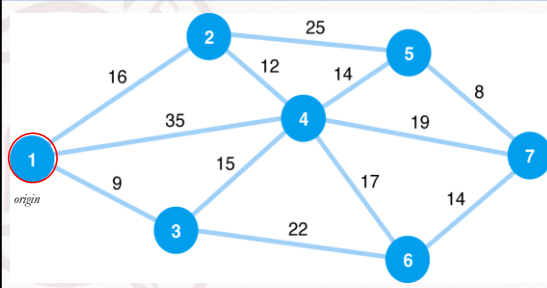
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### The Shortest Route Problem

#### Definition and Example Problem Data (2 of 2)



Source: Garn (2010), Issues in Operations Management Figure 7.3 Network Representation

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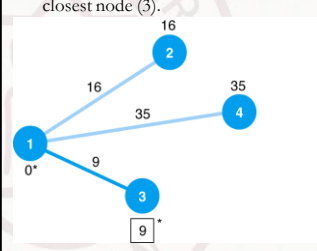
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### The Shortest Route Problem

#### Solution Approach (1 of 8)

Determine the initial shortest route from the origin (node 1) to the closest node (3).



Permanent set	Branch	Time
{1}	1-2	16
	1-4	35
	1-3	9

Figure 7.4 Network with Node 1 in the Permanent Set

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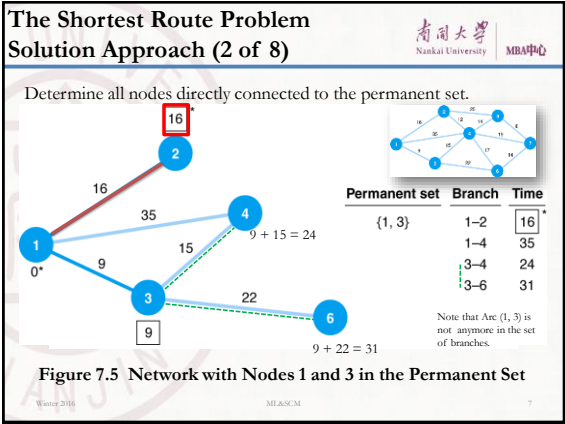
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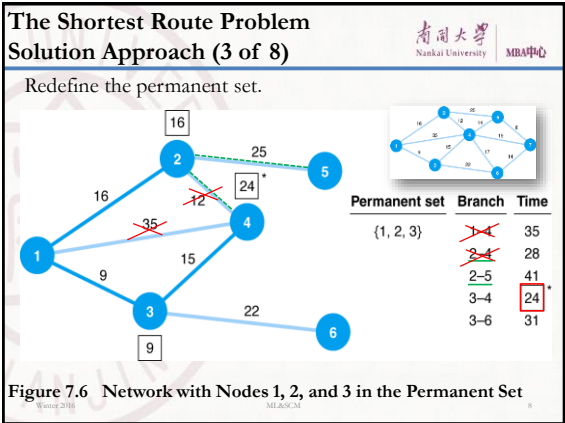
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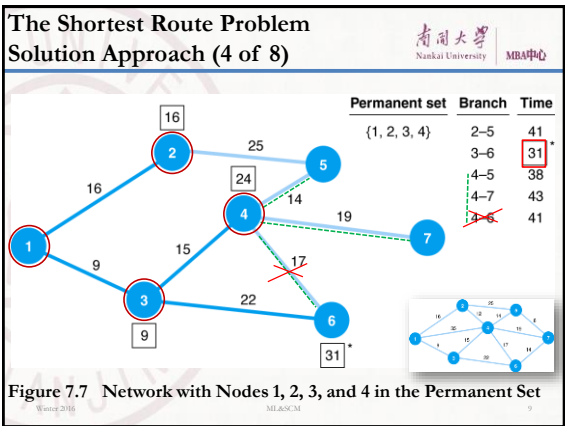
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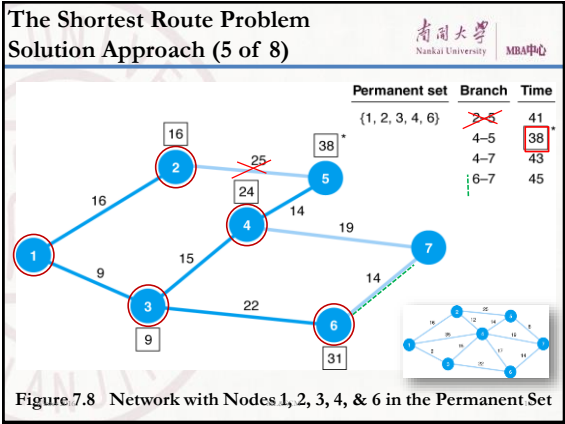
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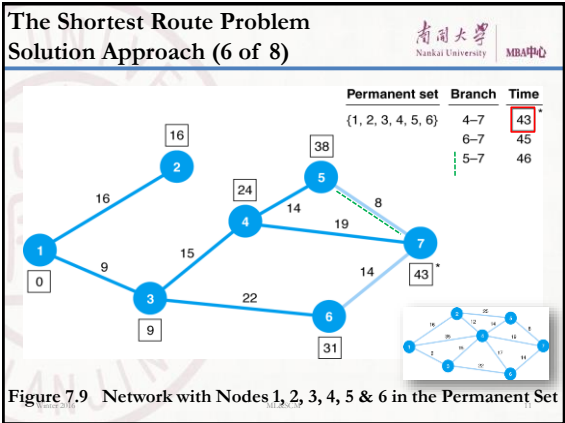
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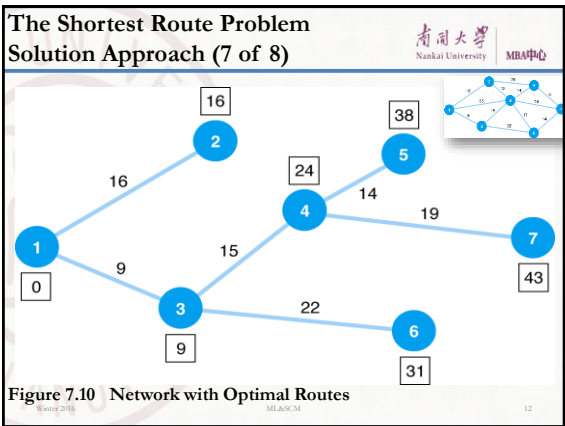
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The Shortest Route Problem

Solution Approach (8 of 8)

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From Los Angeles to:	Route	Total Hours
Salt Lake City (node 2)	1-2	16
Phoenix (node 3)	1-3	9
Denver (node 4)	1-3-4	24
Des Moines (node 5)	1-3-4-5	38
Dallas (node 6)	1-3-6	31
St. Louis (node 7)	1-3-4-7	43

Table 7.1

Shortest Travel Time from Origin to Each Destination

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Shortest Route

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0

min

+

+

↓

↓

↓

↓

↑

↑

↑

↑

1. Put origin in permanent node set and all its outgoing arcs plus "distances" into the branch "analysis" list

2. Look for the minimal distance in this list, this identifies the new arc and node

3. Add the new node to the permanent set

4. Tidy up, i.e. remove arcs that go to the new node

5. Add all its outgoing arcs plus "distances" into the branch "analysis" list

6. Goto step 2 – until all nodes are in the permanent set

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Dijkstra

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Algorithm 1 Dijkstra.

Require:

network  $G = (N, A)$ ; start node  $s$ ; sparse cost matrix  $C = (c_{ij})$

Ensure:

distances  $d$ , predecessors  $p$

1:  $S := \emptyset$ ;  $\bar{S} := N$

2:  $d(i) := \infty$  for each node  $i \in N$ ;

3:  $d(s) := 0$ ;  $p(s) := 0$ ;

4: **while**  $|S| < n$  **do**

5:   let  $i \in \bar{S}$  be a node for which  $d(i) = \min\{d(j) : j \in \bar{S}\}$ ;

6:    $S := S \cup \{i\}$ ;

7:    $\bar{S} := \bar{S} - \{i\}$ ;

8:   **for each**  $(i, j) \in A(i)$  **do**

9:     **if**  $d(j) > d(i) + c_{ij}$  **then**

10:        $d(j) := d(i) + c_{ij}$ ;  $p(j) := i$ ;

11:     **end if**

12:   **end for**

13: **end while**

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### The Shortest Route Problem

#### Integer program

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Formulation as a 0 - 1 integer program (for one destination only)

$x_{ij} = 0$  if branch  $i-j$  is not selected as part of the shortest route and 1 if it is selected.

Minimize  $Z = 16x_{12} + 9x_{13} + 35x_{14} + 12x_{24} + 25x_{25} + 15x_{34} + 22x_{36} + 14x_{45} + 17x_{46} + 19x_{47} + 8x_{57} + 14x_{67}$

subject to:

$x_{12} + x_{13} + x_{14} = 1$

$x_{12} - x_{24} - x_{25} = 0$

$x_{13} - x_{34} - x_{36} = 0$

$x_{14} + x_{24} + x_{34} - x_{45} - x_{46} - x_{47} = 0$

$x_{25} + x_{45} - x_{57} = 0$

$x_{36} + x_{46} - x_{67} = 0$

$x_{47} + x_{57} + x_{67} = 1$

$x_{ij} = 0 \text{ or } 1$

Flow out of source

Conservation-of-flow  
(Note: directed graph)

Flow into destination

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### Application $k$ -Shortest Path

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- Telecommunication – leased lines
- Banks need dedicated line between branches

Austria  
Wolfgang's first Java implementation of a shortest path

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### Bees' tiny brains beat computers

#### – guardian.co.uk

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- “Bees can solve complex mathematical problems which keep computers busy for days, research has shown” – 24<sup>th</sup> October 2010

Dr Nigel Raine, [Royal Holloway, University of London](#)

To be published in: *The American Naturalist*

How many flowers were used?  
•  $10 \Rightarrow 10! = 10 \times 9 \times 8 \times \dots \times 1 = 3,628,800$  possible solutions

How many bees helped solving the problem?  
• Multi Agent Systems, e.g. “Ant algorithms”

Travelling Salesman Problem (TSP)

Dr. A. G. A. J. 18

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The Maximal Flow Problem

Solution Approach (1 of 5)

Step 1: Arbitrarily choose any path through the network from origin to destination and ship as much as possible.

road cars still available

Input

Manchester

2

4

7

2

4

3

6

2

3

5

4

0

4

4

Output

Edinburgh

Rail road cars used

Maximal Flow for Path 1-2-5-6

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The Maximal Flow Problem

Solution Approach (2 of 5)

Step 2: Re-compute branch flow in both directions

Step 3: Select other feasible paths arbitrarily and determine maximum flow along the paths until flow is no longer possible.

8

2

4

0

7

4

3

4

0

3

5

1

4

0

4

0

8

Figure 7.20 Maximal Flow for Path 1-4-6

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The Maximal Flow Problem

Solution Approach (3 of 5)

Continue

14

2

0

7

1

4

3

4

0

3

1

2

6

0

4

0

14

Figure 7.21 Maximal Flow for Path 1-3-6

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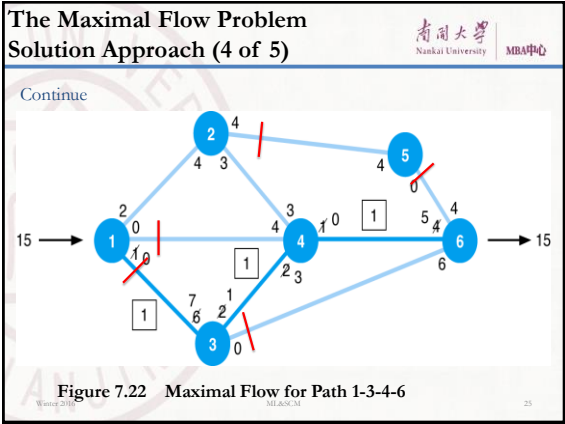
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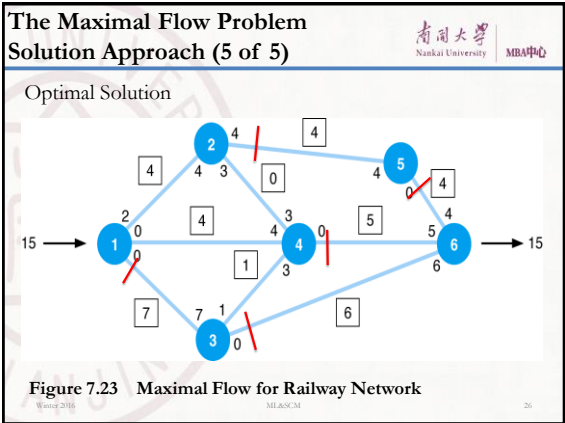
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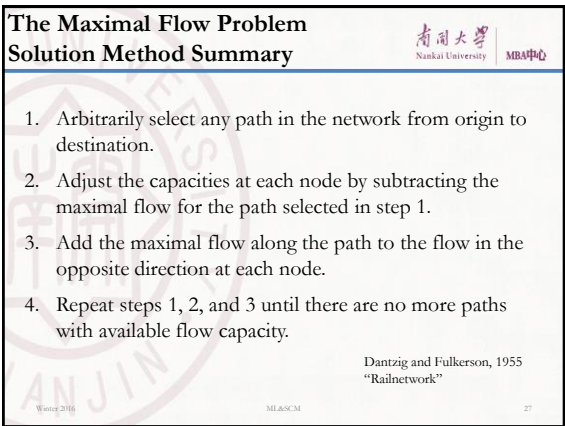
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### The Maximal Flow Problem

#### Integer Program

$x_{ij}$  = flow along branch  $i$ - $j$  and integer

Maximize  $Z = x_{61}$

subject to:

$x_{61} - x_{12} - x_{13} - x_{14} = 0$   
 $x_{12} - x_{24} - x_{25} = 0$   
 $x_{13} - x_{34} - x_{36} = 0$   
 $x_{14} + x_{24} + x_{34} - x_{46} = 0$   
 $x_{25} - x_{56} = 0$

$x_{36} + x_{46} + x_{56} - x_{61} = 0$

Flow in & out of source  
Flow out of source is the same as the flow into the destination  
Conservation of flow (also known as mass balance constraints)

$x_{12} \leq 6 \quad x_{24} \leq 3 \quad x_{34} \leq 2$   
 $x_{13} \leq 7 \quad x_{25} \leq 8 \quad x_{36} \leq 6$   
 $x_{14} \leq 4 \quad x_{46} \leq 5 \quad x_{56} \leq 4$   
 $x_{61} \leq 17 \quad x_{ij} \geq 0 \text{ and integer}$

Note: some opposite flows are missing (for simplification reasons only)

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### General formulation

maximise  $v$  subject to

$$\sum_{\{j:(i,j) \in A\}} x_{ij} - \sum_{\{j:(j,i) \in A\}} x_{ji} = \begin{cases} v & \text{for } i = s \\ 0 & \text{for all } i \in N - \{s \text{ and } t\} \\ -v & \text{for } i = t \end{cases}$$

$0 \leq x_{ij} \leq u_{ij}$  for each  $(i, j) \in A$

$x = \{x_{ij}\}$  flow,  $v$  value of flow

Assumptions:

- The network is directed
- All capacities are nonnegative integers
- There is a path between  $s$  and  $t$  with finite capacity

More assumptions:

- No directed path from  $s$  to  $t$  entirely with infinite capacities
- If  $(i, j) \in A$  then  $(j, i) \in A$
- No parallel arcs

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### Applications

- Flow of water, electricity, gas and oil
  - E.g. provide required gas to London using all available supply routes

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### Applications

- Traffic flow, Railway system
  - Transport as many goods as possible with time being constraint
  - Joint car navigation system
- Flow of Products through a production line, conveyor systems or sorting centre
  - Maximise letters sorted using all available paths



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### Applications

- Tanker Scheduling (Dantzig and Fulkerson, 1954)
- Scheduling of parallel Machines (Federgruen, 1986)
- Distributed Computing (Stone, 1977)



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### Recap

- What is the shortest path?
  - A path from a source node to a destination node that minimises the flow.
- What is the maximal flow?
  - A network utilisation that maximises the flow from a source node to a destination node.
- What is probably the most popular shortest path commercial product?
  - SatNav

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### The End

- Any questions?



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### Appendix

- Another shortest route example
- Shortest route in Excel
- Maximal flow in Excel

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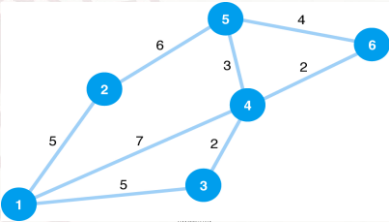
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### Problem Statement and Data

- Determine the shortest route from Atlanta (node 1) to each of the other five nodes (branches show travel time between nodes).
- Assume branches show distance (instead of travel time) between nodes, develop a minimal spanning tree.



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Shortest Route Solution (1 of 2)

Step 1 (part A): Determine the Shortest Route Solution

1.	Permanent Set	Branch	Time
	{1}	1-2	[5]
		1-3	5
		1-4	7
2.	{1,2}	1-3	[5]
		1-4	7
		2-5	11
3.	{1,2,3}	1-4	[7]
		2-5	11
		3-4	7
4.	{1,2,3,4}	4-5	10
		4-6	[9]
5.	{1,2,3,4,6}	4-5	[10]
		6-5	13
6.	{1,2,3,4,5,6}		

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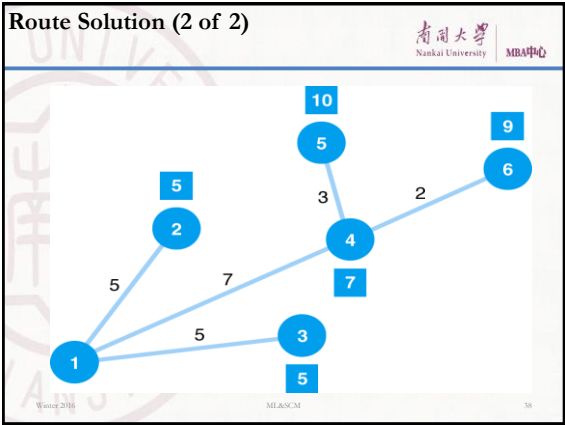
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The Shortest Route Problem

Integer program

Formulation as a 0 - 1 integer linear programming problem.

$x_{ij} = 0$  if branch  $i-j$  is not selected as part of the shortest route and 1 if it is selected.

Minimize  $Z = 16x_{12} + 9x_{13} + 35x_{14} + 12x_{24} + 25x_{25} + 15x_{34} + 22x_{36} + 14x_{45} + 17x_{46} + 19x_{47} + 8x_{57} + 14x_{67}$

subject to:

- $x_{12} + x_{13} + x_{14} = 1$  → Flow out of source
- $x_{12} - x_{24} - x_{25} = 0$
- $x_{13} - x_{34} - x_{36} = 0$
- $x_{14} + x_{24} + x_{34} - x_{45} - x_{46} - x_{47} = 0$  → Conservation-of-flow
- $x_{25} + x_{45} - x_{57} = 0$
- $x_{36} + x_{46} - x_{67} = 0$
- $x_{47} + x_{57} + x_{67} = 1$  → Flow into destination
- $x_{ij} = 0 \text{ or } 1$

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The Shortest Route Problem

Computer Solution with Excel (2 of 4)

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Total hours

Select Branch	City	Distance
1	Los Angeles	16
2	Los Angeles	35
3	Los Angeles	35
4	Los Angeles	35
5	Los Angeles	35
6	Los Angeles	35
7	Los Angeles	35
8	Los Angeles	35
9	Los Angeles	35
10	Los Angeles	35
11	Los Angeles	35
12	Los Angeles	35
13	Los Angeles	35
14	Los Angeles	35
15	Los Angeles	35
16	Los Angeles	35
17	Los Angeles	35
18	Los Angeles	35
19	Los Angeles	35
20	Los Angeles	35

First constraint;  
=A6+A7+A8

Constraint for node 2;  
=A6-A9-A10

Decision variables

Exhibit 7.3

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The Shortest Route Problem

Computer Solution with Excel (3 of 4)

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One truck leaves node 1, and one truck ends at node 7.

Flow constraints

Solver Parameters

Set Target Cell: \$B\$11 To: Value of: 0

By Changing Variable Cells: \$A\$6:\$A\$17

Subject to the Constraints: \$A\$6:\$A\$17 = integer \$B\$12 = 1 \$B\$7:\$B\$11 = 0

Exhibit 7.4

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The Shortest Route Problem

Computer Solution with Excel (4 of 4)

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One truck flows out of node 1; one truck flows into node 7.

Select Branch	City	Distance
1	Los Angeles	16
2	Los Angeles	35
3	Los Angeles	35
4	Los Angeles	35
5	Los Angeles	35
6	Los Angeles	35
7	Los Angeles	35
8	Los Angeles	35
9	Los Angeles	35
10	Los Angeles	35
11	Los Angeles	35
12	Los Angeles	35
13	Los Angeles	35
14	Los Angeles	35
15	Los Angeles	35
16	Los Angeles	35
17	Los Angeles	35
18	Los Angeles	35
19	Los Angeles	35
20	Los Angeles	35

Exhibit 7.5

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The Maximal Flow Problem

Computer Solution with Excel (1 of 4)

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$x_{ij}$  = flow along branch  $i$ - $j$  and integer

Maximize  $Z = x_{61}$

subject to:

$x_{61} - x_{12} - x_{13} - x_{14} = 0$

$x_{12} - x_{24} - x_{25} = 0$

$x_{13} - x_{34} - x_{36} = 0$

$x_{14} + x_{24} + x_{34} - x_{46} = 0$

$x_{25} - x_{56} = 0$

$x_{36} + x_{46} + x_{56} - x_{61} = 0$

$x_{12} \leq 6 \quad x_{24} \leq 3 \quad x_{34} \leq 2$

$x_{13} \leq 7 \quad x_{25} \leq 8 \quad x_{36} \leq 6$

$x_{14} \leq 4 \quad x_{46} \leq 5 \quad x_{56} \leq 4$

$x_{61} \leq 17 \quad x_{ij} \geq 0 \text{ and integer}$

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The Maximal Flow Problem

Computer Solution with Excel (2 of 4)

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Objective—maximize flow from node 6

Constraint at node 1;  $=C15-C6-C7-C8$

Constraint at node 6;  $=C12+C13+C14-C15$

Decision variables

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Exhibit 7.8

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The Maximal Flow Problem

Computer Solution with Excel (3 of 4)

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Flow into and out of nodes must equal each other.

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Exhibit 7.9

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The Maximal Flow Problem

Computer Solution with Excel (4 of 4)

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Exhibit 7.10

Branch	From Node	To Node	Flow	Capacity
1	1	2	4	6
2	1	3	7	7
3	1	4	4	4
4	2	4	0	3
5	2	5	4	6
6	3	4	1	2
7	3	6	6	6
8	4	6	5	5
9	5	6	4	4
10	6	1	15	17
Total			15	

Node	Network Flow
1	0
2	0
3	0
4	0
5	0
6	0

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The Shortest Route Problem

Solution Method Summary

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1. Select the node with the shortest direct route from the origin.
2. Establish a permanent set with the origin node and the node that was selected in step 1.
3. Determine all nodes directly connected to the permanent set of nodes.
4. Select the node with the shortest route from the group of nodes directly connected to the permanent set of nodes.
5. Repeat steps 3 & 4 until all nodes have joined the permanent set.

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Shortest Route - Algorithm

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INPUT: nodes  $N$ , Arcs  $A$ , Costs  $C$

OUTPUT: shortest paths (predecessor notation)

Dijkstra, 1959

$S := \emptyset, \bar{S} := N$   
 $\forall i \in N: p_i := 0, d_i := \infty$  set predecessors and node distances  
 $d_s := 0$  distance from source  
while  $|S| < n$  until all paths are found  
     $i \in \bar{S}: d_i := \min_{j \in S} d_j$  find node with minimal distance  
     $S := S \cup i, \bar{S} := \bar{S} \setminus i$  swap element between sets  
     $\forall (i, j) \in A_i$ : check all arcs going out from  $i$   
         $d_j > d_i + c_{ij} \Rightarrow$  if distance can be reduced  
         $d_j := d_i + c_{ij}, p_j := p_i$  then set distance and predecessor

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